

Formalizing the Wu-Ritt Characteristic Set Method in Lean 4

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Abstract. We formalize the Wu-Ritt characteristic set method in Lean 4 for triangular decomposition of multivariate polynomial systems. The development covers polynomial orders, initials, pseudo-division, triangular and ascending sets, basic and characteristic sets, and algorithms for characteristic-set computation and zero decomposition, with termination and correctness proofs. In particular, we verify the well-ordering principle and the zero decomposition theorem, providing a machine-checked foundation for certified polynomial system solving and geometry theorem proving.

Keywords: Wu-Ritt characteristic set method, ascending sets, pseudo-division, characteristic sets, Lean

1 Introduction

The Wu-Ritt characteristic set method⁴ introduced by Wu Wen-Tsun [6], is a fundamental tool for solving systems of multivariate polynomial equations and for automated theorem proving in elementary geometry. In this paper, we present a formalization of the Wu-Ritt characteristic set method in Lean 4 [3], following Wu’s mathematical and algorithmic treatment [7]. Our development formalizes the basic concepts of the theory, including pseudo-division, pseudo-remainder with respect to a polynomial or a triangular set, ascending sets, basic sets, and characteristic sets. We verify not only the main algorithms of the method but also the central statements of the Wu-Ritt theory. In particular, we formalize the well-ordering principle that connects a polynomial system with its characteristic set, and we prove the correctness of the zero decomposition algorithm, showing that it expresses the zero set of the original system as a finite union of zero sets of triangular sets, away from the zeros of the corresponding initials.

In Coq [4], a formalization of Wu’s simple method was previously carried out for proving geometric theorems [1]. Our work provides a comprehensive formalization of the characteristic set method in Lean 4. It builds on Mathlib’s algebraic infrastructure and provides a basis for certified symbolic computation and automated geometry theorem proving in the Lean ecosystem [5].

2 Preliminaries

Let R be a commutative ring with identity, and let σ be a finite set of indices equipped with a linear order $<$. We write $R[X_i]_{i \in \sigma}$ for the polynomial ring under consideration, ordering the variables by declaring

$$X_i < X_j \quad \text{whenever } i < j.$$

⁴ : The terminology “Wu-Ritt” is used here to emphasize Wu Wen-Tsun’s algorithmic development of the characteristic set method, built on Ritt’s earlier theory of characteristic sets. In particular, we follow Wu’s terminology and treatment in [6].

Let $p \in R[X_i]_{i \in \sigma}$ be a polynomial. We denote the greatest variable appearing in p by X_{m_p} and call it the *main variable* of p , where $m_p \in \sigma$ is the index of that variable. We define the *main degree* of p , denoted by $\deg(p)$, to be the degree of p with respect to its main variable, namely $\deg_{m_p}(p)$. If p is a constant polynomial, then its main variable is undefined, and we define $\deg(p) = 0$.

Definition 1 (Polynomial Order). *Let p and q be two non-zero multivariate polynomials with main variables X_{m_p}, X_{m_q} , respectively. We say p is less than q , denoted by $p < q$, if either $m_p < m_q$ or $m_p = m_q$ and $\deg(p) < \deg(q)$.*

The order $<$ can be viewed as the lexicographic order on $(m_p, \deg(p))$, comparing first the main-variable index and then the degree. Unlike Gröbner-basis monomial orders, which compare monomials to define leading terms and reductions [2], it compares whole polynomials for our triangular/pseudo-division procedure. Thus, it is neither induced by nor intended to extend such monomial orders; apart from `MvPolynomial`, the two formalizations are largely independent.

```
def order (p : MvPolynomial σ R) : WithBot σ ×l ℕ := (p.vars.max, p.mainDegree)
```

We say that p and q have the same ordering, denoted by $p \sim q$, if neither $p < q$ nor $q < p$ holds. Equivalently,

$$p \sim q \iff m_p = m_q \wedge \deg(p) = \deg(q).$$

We write $p \preceq q$ if $p < q$ or $p \sim q$. The relation $\preceq$ defines a preorder on $R[X_i]_{i \in \sigma}$, and $\sim$ is the corresponding equivalence relation. Hence $\preceq$ induces a partial order on the quotient $R[X_i]_{i \in \sigma} / \sim$.

```
instance : Preorder (MvPolynomial σ R) where
  le := InvImage (· ≤ ·) order
  le_refl := fun _ => by rw [InvImage]
  le_trans := fun _ _ => le_trans

theorem lt_def : p < q ↔ p.vars.max < q.vars.max ∨
  p.vars.max = q.vars.max ∧ p.mainDegree < q.mainDegree

instance : Setoid (MvPolynomial σ R) :=
  AntisymmRel.setoid (MvPolynomial σ R) (· ≤ ·)
```

Definition 2 (Initial). *Let p be a non-constant polynomial with main variable X_i . Writing p as a polynomial in X_i , $p = c_d X_i^d + c_{d-1} X_i^{d-1} + \dots + c_0$, where each $c_k \in R[X_j]_{j < i}$ and $c_d \neq 0$, we call c_d the *initial* of p , and denote it by $\text{init}(p)$ or simply I_p . We define the initial of the zero polynomial to be 0. If p is a nonzero constant, then we set $I_p = 1$.*

```
def initialOf (p : MvPolynomial σ R) (i : σ) : MvPolynomial σ R :=
  Σ s ∈ p.support with s i = p.degreeOf i, monomial (s.erase i) (p.coeff s)

def initial (p : MvPolynomial σ R) : MvPolynomial σ R :=
  if p = 0 then 0 else
    match p.vars.max with
    | ⊥ => 1
    | some m => p.initialOf m
```

```
theorem initialOf_eq_leadingCoeff [DecidableEq σ] {p : MvPolynomial σ R} {i : σ} :
  p.initialOf i = rename Subtype.val
    (optionEquivLeft R {b // b ≠ i} (rename (Equiv.optionSubtypeNe i).symm p)).leadingCoeff
```

The *initial product* of polynomials $p_1, p_2, \dots, p_r$ is defined to be the product of their initials $\prod_{i=1}^r I_{p_i}$.

Example 1. Consider $p = 5xy^2 + x^3y, q = y^2 + x^2y, r = x^5 \in \mathbb{Z}[x, y]$ with $x < y$. Then $r < p$ and $p \sim q$, $I_p = \text{init}_y(p) = 5x, I_q = 1$, and $I_r = 1$.

Definition 3 (Reduce). *Let p be a non-constant polynomial with main variable X_{m_p} . A polynomial q is said to be reduced with respect to p if the degree of q in X_{m_p} is strictly smaller than the main degree of p . By convention, the zero polynomial is reduced with respect to every polynomial. If p is a constant polynomial, then no nonzero polynomial is reduced with respect to p .*

```
def reducedTo (q p : MvPolynomial σ R) : Prop :=
  if q = 0 then True
  else
    match p.vars.max with
    | ⊥ => False
    | some m => q.degreeOf m < p.degreeOf m
```

Let PS be a set of polynomials. We say that a polynomial q is *reduced with respect to PS* if it is reduced with respect to every polynomial in PS .

3 Triangular Set

Definition 4 (Triangular Set). A set of polynomials S is called a triangular set if all its elements are nonzero and can be arranged as a sequence

$$s_1, s_2, \dots, s_r \in R[X_i]_{i \in \sigma},$$

whose main variables $X_{m_1}, X_{m_2}, \dots, X_{m_r}$ satisfy

$$m_1 < m_2 < \dots < m_r.$$

```
structure TriangularSet (σ R : Type*) [CommSemiring R] [LinearOrder σ] where
  /-- The underlying list of non-zero polynomials -/
  toList : List (MvPolynomial σ R)
  /-- No polynomial in the list is zero -/
  zero_notMem' : 0 ∉ toList
  /-- The main variables of successive polynomials are strictly increasing -/
  isChain_max_vars' : toList.IsChain (·.vars.max < ·.vars.max)

instance : FunLike (TriangularSet σ R) ℕ (MvPolynomial σ R) where
  coe S n := S.toList[n]?.getD 0
  coe_injective' := ...

def length (S : TriangularSet σ R) : ℕ := S.toList.length
```

Especially, our formalization allows s_1 to be a nonzero constant. In this case, m_1 is undefined, while $m_2 < \dots < m_r$.

When S is a triangular set, we always write it as $s_1, \dots, s_r$ so that the main variables of the s_i are strictly increasing with i .

Definition 5 (Ordering of Triangular Sets). Let $S = (s_1, \dots, s_r)$ and $S' = (s'_1, \dots, s'_{r'})$ be two triangular sets. We say that S is less than S' , and write $S \prec S'$, if one of the following conditions holds:

1. there exists $1 \leq k \leq \min(r, r')$ such that

$$s_1 \sim s'_1, s_2 \sim s'_2, \dots, s_{k-1} \sim s'_{k-1}, \text{ while } s_k \prec s'_k.$$

2. $r > r'$ and $s_i \sim s'_i$ for all $1 \leq i \leq r'$.

```
-- The indices here are 0-based.
def order (S : TriangularSet σ R) : Lex (ℕ → WithTop (WithBot σ ×l ℕ)) :=
  fun i ↦ if i < S.length then WithTop.some (S i).order else ⊤
```

```
theorem lt_def : S < T ↔ (∃ k < S.length, S k < T k ∧ ∀ i < k, S i ≈ T i) ∨
  (T.length < S.length ∧ ∀ i < T.length, S i ≈ T i)
```

As in Definition 1, we extend the relations $\preceq$ and $\sim$ to triangular sets. The relation $\preceq$ is a preorder on triangular sets, and it induces a partial order on the quotient $\text{TriangularSet}/\sim$.

The polynomial order is well-founded whenever σ is well-ordered. For triangular sets, however, well-foundedness requires the stronger assumption that σ is **finite**.

Proposition 1 (Well-ordering property of triangular sets) *If σ is finite, then there is no infinite strictly decreasing sequence of triangular sets.*

```
instance [Finite  $\sigma$ ] : WellFoundedLT (TriangularSet  $\sigma$  R)
```

This well-ordering property is a key ingredient in proving the termination of the algorithms for computing characteristic sets and zero decompositions.

Definition 6 (Ascending Set). A triangular set $AS = \{as_1, \dots, as_r\}$ is called an ascending set if, for every pair i, j with $i < j$, the polynomial as_j is reduced with respect to as_i .

```
class AscendingSetTheory ( $\sigma$  R : Type*) [CommSemiring R] [DecidableEq R] [LinearOrder  $\sigma$ ] where
  /-- The reduction relation used to define the ascending property. -/
  protected reducedTo' : MvPolynomial  $\sigma$  R → MvPolynomial  $\sigma$  R → Prop
  decidableReducedTo : DecidableRel reducedTo' := by infer_instance
  /-- A key property linking the ascending set structure to the initial.
  If 'S' is an ascending set, the initial of any non-constant element in 'S'
  must be reduced with respect to 'S'. -/
  protected initial_reducedToSet_of_max_vars_ne_bot :  $\forall$  {S : TriangularSet  $\sigma$  R} {i :  $\mathbb{N}$ },
    ( $\forall$  {j}, i < j → j < S.length → reducedTo' (S j) (S i)) →
    (S i).vars.max  $\neq$   $\perp$  → (S i).initial.reducedToSet S

def isAscendingSet (S : TriangularSet  $\sigma$  R) : Prop :=
   $\forall$  {i j}, i < j → j < S.length → AscendingSetTheory.reducedTo' (S j) (S i)
```

The essential property of an ascending set AS is that, for every non-constant polynomial $p \in AS$, its initial I_p is reduced with respect to AS itself. This property plays a crucial role in the termination proof of the zero decomposition algorithm.

Definition 7 (Basic Set). A basic set BS of a polynomial set PS is the minimal ascending set contained in PS .

```
class HasBasicSet ( $\sigma$  R : Type*) [CommSemiring R] [DecidableEq R] [LinearOrder  $\sigma$ ]
  extends AscendingSetTheory  $\sigma$  R where
  /-- Computes a Basic Set from a list of polynomials. -/
  basicSet : List (MvPolynomial  $\sigma$  R) → TriangularSet  $\sigma$  R
  basicSet_isAscendingSet (l : List (MvPolynomial  $\sigma$  R)) : (basicSet l).isAscendingSet
  basicSet_subset (l : List (MvPolynomial  $\sigma$  R)) :  $\forall$  {c}, c  $\in$  basicSet l → c  $\in$  l
  basicSet_minimal (l : List (MvPolynomial  $\sigma$  R)) :
     $\forall$  {S}, S.isAscendingSet → ( $\forall$  {p}, p  $\in$  S → p  $\in$  l) → basicSet l  $\leq$  S
  /-- Order reduction property: appending a reduced element strictly decreases the basic set order.
  Crucial for proving termination of zero decomposition. -/
  basicSet_append_lt_of_exists_reducedToSet :  $\forall$  {l1 l2 : List (MvPolynomial  $\sigma$  R)},
    ( $\exists$  p  $\in$  l2, p  $\neq$  0  $\wedge$  p.reducedToSet (basicSet l1)) → basicSet (l2 ++ l1) < basicSet l1
```

The next proposition is used to show that, throughout the algorithm, the order of the basic set strictly decreases.

Proposition 2 Let p be a polynomial reduced with respect to the basic set of PS . Then the basic set of $PS \cup \{p\}$ is strictly smaller than that of PS under the prescribed ordering.

```
theorem basicSet_append_lt_of_exists_reducedToSet
  (h :  $\exists$  p  $\in$  l2, p  $\neq$  0  $\wedge$  p.reducedToSet l1.basicSet) : (l2 ++ l1).basicSet < l1.basicSet
```

4 Polynomial Pseudo-Division

Theorem 3 (Pseudo-Remainder) Let $f, g \in R[X_i]_{i \in \sigma}$, and suppose that f is a non-constant polynomial with initial $I = \text{init}(f)$. Then there exist $s \in \mathbb{N}$, $q \in R[X_i]_{i \in \sigma}$, and $r \in R[X_i]_{i \in \sigma}$ such that

$$I^s \cdot g = q \cdot f + r,$$

where r is reduced with respect to f . The polynomial r is called the pseudo-remainder of g with respect to f , denoted by

$$r = \text{Remdr}(g/f).$$

The following procedure gives an algorithm for computing s , q , and r .

Require: polynomial g , nonzero polynomial f with main variable X_m

Ensure: pseudo-remainder r of g divided by f

```

1:  $s := 0, q := 0, r := g$ 
2:  $I := \text{init}_m(f)$ 
3:  $d_f := \deg_m(f)$ 
4:  $d_r := \deg_m(r)$ 
5: while  $d_r \geq d_f$  do
6:    $t := \text{init}_m(r) \cdot X_m^{d_r - d_f}$ 
7:    $r := I \cdot r - t \cdot f$ 
8:    $q := I \cdot q + t$ 
9:    $s := s + 1$ 
10:   $d_r := \deg_m(r)$ 
11: end while
12: return  $r$ 

```

The key step is Step 6, where r is multiplied by I , while f is multiplied by $\text{init}_m(r)$, so that the leading term of r is canceled. The factor $X_m^{d_r - d_f}$ aligns the degrees of r and f in the main variable X_m . Consequently, the degree of r strictly decreases at each iteration.

The general `pseudo` function is defined over a field: it automatically detects the main variable and also handles constant divisors. Hence, from this point on, we assume that R is a field.

```

variable {R : Type*} [Field R] [DecidableEq R] [LinearOrder σ] (g f : MvPolynomial σ R)

def pseudo : PseudoResult (MvPolynomial σ R) :=
  if f = 0 then ⟨0, 0, g⟩
  else
    match f.vars.max with
    | ⊥ => ⟨0, (f.coeff 0)-1 · g, 0⟩
    -- pseudoOf implements the pseudo-division algorithm described above
    | some m => g.pseudoOf m f

theorem pseudo_equation : f.initial ^ (g.pseudo f).exponent * g = (g.pseudo f).quotient * f + (g.pseudo
  f).remainder

```

This allows us to define the pseudo-remainder with respect to a triangular set.

Theorem 4 (Pseudo-remainder with respect to a triangular set) *Let $S = (s_1, \dots, s_k)$ be a triangular set and let $g \in R[X_i]_{i \in \sigma}$. Then there exist*

$$r, q_1, \dots, q_k \in R[X_i]_{i \in \sigma}$$

and $e_1, \dots, e_k \in \mathbb{N}$ such that

$$\left(\prod_{i=1}^k \text{init}(s_i)^{e_i} \right) g = \sum_{i=1}^k q_i s_i + r,$$

where r is reduced with respect to S . The polynomial r is called the pseudo-remainder of g with respect to S and is denoted by

$$r = \text{Remdr}(g/S).$$

Proof. This remainder can be computed by recursively performing pseudo-division, starting from the last element of S . First, we perform pseudo-division of g by s_k , obtaining e_k, q_k , and r_k . We then perform pseudo-division of r_k by s_{k-1} , obtaining e_{k-1}, q_{k-1} , and r_{k-1} . Continuing in this way, we eventually obtain a final remainder, r_0 . This polynomial r_0 is the pseudo-remainder of g with respect to S .

```

def setPseudo.go (f : ℕ → MvPolynomial σ R) (fuel : ℕ) (es : List ℕ)
  (qs : List (MvPolynomial σ R)) (r : MvPolynomial σ R) : SetPseudoResult (MvPolynomial σ R) :=
  if fuel = 0 then ⟨es, qs, r⟩
  else
    let p := r.pseudo (f (fuel - 1))
    let es' := p.exponent :: es

```

```

let qs' := p.quotient :: qs.map (· * (f (fuel - 1)).initial ^ p.exponent)
let r' := p.reminder
go f (fuel - 1) es' qs' r'

def setPseudo (S : TriangularSet σ R) : SetPseudoResult (MvPolynomial σ R) :=
  setPseudo.go S S.length [] [] g

theorem setPseudo_equation : letI result := g.setPseudo S
  (∀ i : Fin result.exponents.length, (S i).initial ^ result.exponents[i]) * g
  = (Σ i : Fin result.quotients.length, result.quotients[i] * S i) + result.reminder

```

To decouple the algorithm from the correctness proofs, the notion of remainder is defined as a proposition, not tied to a specific implementation:

```

def IsRemainder (r g f : MvPolynomial σ R) : Prop :=
  r.reducedTo f ∧ ∃ (s : ℕ) (q : MvPolynomial σ R), f.initial ^ s * g = q * f + r

def IsSetRemainder (r g : MvPolynomial σ R) (S : TriangularSet σ R) : Prop := r.reducedToSet S ∧
  ∃ (es : List ℕ) (qs : List (MvPolynomial σ R)), (es.length = S.length ∧ qs.length = S.length) ∧
  (∀ i : Fin es.length, (S i).initial ^ es[i] * g = (Σ i : Fin qs.length, qs[i] * S i) + r

```

5 Characteristic Set

For a polynomial set PS , we write $\text{Zero}(PS)$ for its vanishing set, that is, the set of all points at which every polynomial in PS evaluates to zero. We also use the abbreviation

$$\text{Zero}(CS/IP) := \text{Zero}(CS) \setminus \text{Zero}(IP),$$

namely the set of common zeros of CS at which the initial product IP does not vanish. In Lean, these are defined as $MvPolynomial.vanishingSet$ and $MvPolynomial.vanishingSet'$, which coincide with the zero locus of the ideal generated by the set.

Definition 8 (Characteristic Set). *A triangular set CS is a characteristic set of PS if $\forall p \in PS, \text{Remdr}(p/CS) = 0$ and $\text{Zero}(PS) \subseteq \text{Zero}(CS)$.*

```

def TriangularSet.IsCharacteristicSet {α : Type*} [Membership (MvPolynomial σ R) α] (K : Type*)
  [CommSemiring K] [Algebra R K] (CS : TriangularSet σ R) (a : α) : Prop :=
  (∀ g ∈ a, (0 : MvPolynomial σ R).IsSetRemainder g CS) ∧ vanishingSet K a ⊆ vanishingSet K CS

```

Theorem 5 (Well-ordering Principle) *Let $CS = (c_1, \dots, c_r)$ be a characteristic set of PS , and let $IP = \prod_{i=1}^r I_{c_i}$ be the initial product of CS . Then $\text{Zero}(CS/IP) \subseteq \text{Zero}(PS)$, and consequently $\text{Zero}(CS/IP) = \text{Zero}(PS/IP)$. Moreover,*

$$\text{Zero}(PS) = \text{Zero}(CS/IP) \cup \left(\bigcup_{p \in CS} \text{Zero}(PS \cup \{I_p\}) \right). \quad (1)$$

```

theorem vanishingSet_diff_initialProd_subset
  (h : (∀ g ∈ PS, (0 : MvPolynomial σ R).IsSetRemainder g CS)) :
  vanishingSet K CS \ vanishingSet' K (initialProd CS.toFinset) ⊆ vanishingSet K PS

theorem vanishingSet_diff_initialProd_eq (h : CS.IsCharacteristicSet K PS) :
  vanishingSet K CS \ vanishingSet' K (initialProd CS.toFinset) =
  vanishingSet K PS \ vanishingSet' K (initialProd CS.toFinset)

theorem vanishingSet_decomposition (h : CS.IsCharacteristicSet K PS) : vanishingSet K PS =
  vanishingSet K CS \ vanishingSet' K (initialProd CS.toFinset) ∪
  (∪ p ∈ CS, vanishingSet K PS ∩ vanishingSet' K p.initial)

```

Theorem 6 *A characteristic set CS can be constructed in a finite number of steps using the following algorithm:*

Require: $PS = (p_1, \dots, p_s)$
Ensure: a Characteristic set $CS = (c_1, \dots, c_t)$ for PS ,
1: $PS' := PS$
2: **repeat**
3: $RS := \emptyset$
4: $BS := \text{BasicSet}(PS')$
5: **for** each $\{p\}$ in PS' **do**
6: $r := \text{Remdr}(p, BS)$
7: **if** $r \neq 0$ **then**
8: $RS := RS \cup \{r\}$
9: **end if**
10: **end for**
11: $PS' := PS \cup RS \cup BS$
12: **until** $RS = \emptyset$
13: **return** $CS := BS$

```
def characteristicSet.go [Finite σ] (l₀ l : List (MvPolynomial σ R))
  : TriangularSet σ R :=
  let BS := l.basicSet.val
  let lBS := BS.toList
  let RS : List _ := ((l \ lBS).map fun p ↦ (p.setPseudo BS).remainder).filter (· ≠ 0)
  if RS = [] then BS
  else go l₀ (l₀ ++ RS ++ lBS)
  termination_by l.basicSet

def characteristicSet [Finite σ] (l) : TriangularSet σ R := characteristicSet.go l l
```

Theorem 7 (Zero Decomposition Theorem) *For a polynomial system PS , there exists a finite collection of triangular sets*

$$\mathcal{ZD} = \{CS_j\}_j$$

such that

$$\text{Zero}(PS) = \bigcup_j \text{Zero}(CS_j/IP_j), \quad (2)$$

where IP_j denotes the initial product of CS_j . Moreover, for every j and every $p \in PS$,

$$\text{Remdr}(p/CS_j) = 0.$$

```
def zeroDecomposition [Finite σ] (l : List (MvPolynomial σ R)) : List (TriangularSet σ R) :=
  let CS := l.characteristicSet
  let subDecomp := (CS.toList.filter fun p ↦ p.vars.max ≠ 1).attach.map
    fun ⟨p, _⟩ ↦ zeroDecomposition (p.initial :: CS.toList ++ l)
  CS :: subDecomp.flatten
  termination_by l.basicSet
  decreasing_by ...
```

```
theorem vanishingSet_eq_zeroDecomposition_union :
  vanishingSet K l = ⋃ CS ∈ l.zeroDecomposition,
    vanishingSet K CS \ vanishingSet' K (initialProd CS.toFinset)

theorem zero_isSetRemainder_of_mem_zeroDecomposition :
  ∀ CS ∈ l.zeroDecomposition, ∀ g ∈ l, (0 : MvPolynomial σ R).IsSetRemainder g CS
```

For brevity, we omit the proof of Theorem 7 and provide an example illustrating the zero-decomposition procedure.

Example 2. Consider $PS = \{x^2 + y^2 - 1, xy\}$ over the rational number field, where $x < y$. The characteristic set of PS can be taken as $CS = \{x^3 - x, xy\}$. The initials of the two polynomials in CS are 1 and x , so the only nontrivial initial is x . Thus (1) gives

$$\text{Zero}(PS) = \text{Zero}(CS/\{x\}) \cup \text{Zero}(PS \cup \{x\}).$$

Since $x \neq 0$ on the first branch, $xy = 0$ gives $y = 0$, and $x^3 - x = 0$ gives $x^2 - 1 = 0$. Hence $\text{Zero}(CS/\{x\}) = \text{Zero}(x^2 - 1, y)$. On the second branch, $PS \cup \{x\}$ has characteristic set $\{x, y^2 - 1\}$, whose initials are all 1, so the recursion stops. Therefore

$$\text{Zero}(PS) = \text{Zero}(x^2 - 1, y) \cup \text{Zero}(x, y^2 - 1) = \{(\pm 1, 0), (0, \pm 1)\}.$$

6 Conclusions

We formalized the Wu-Ritt characteristic set method in Lean 4, including polynomial orders, pseudo-division, triangular and ascending sets, characteristic sets, and zero decomposition. We proved termination and correctness, including the well-ordering principle and the zero-set decomposition theorem.

The development consists of about 3000 lines of `sorry`-free Lean 4 code based on Mathlib 4 (v4.29.0), and is available under the Apache-2.0 license at https://github.com/WuProver/lean_characteristic_set. Since `MvPolynomial` is noncomputable, our current focus is on correctness proofs. Future work includes executable extraction and extensions to differential polynomials.

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